

## The Nature of Design

The design problem is then reduced to finding the inverse in order to produce x. This is what is here called *ideal* design

 $\longrightarrow \mathbf{f}^{-1}(*) \longrightarrow \mathbf{f}(*)$ 

In the general case, however, the inverse can not be found analytically, since is nonlinear in the general case.

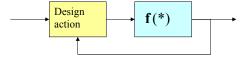
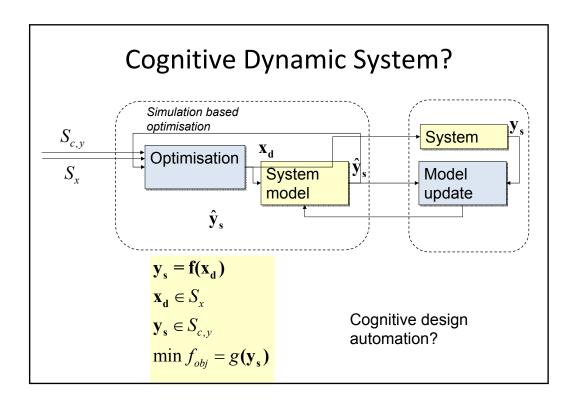
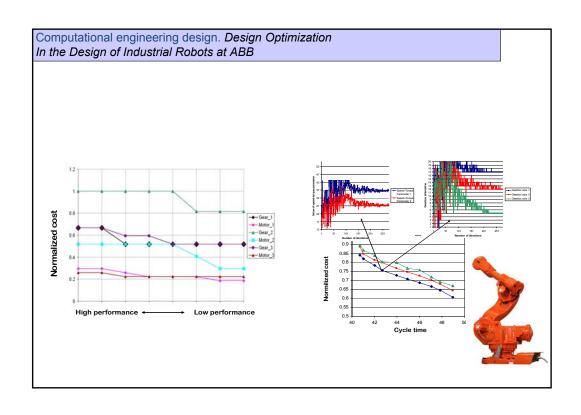
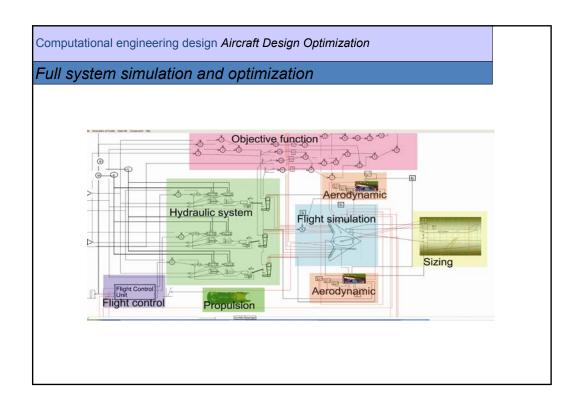
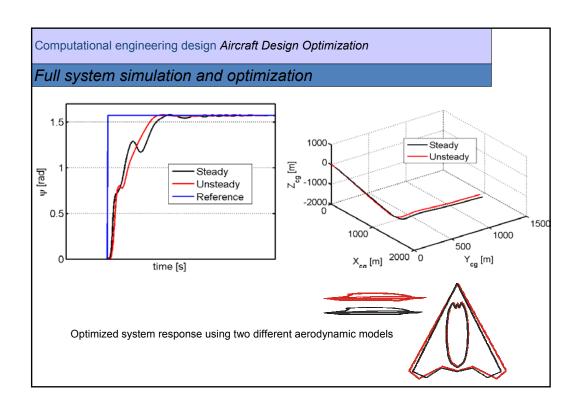


Figure 2. Iterative design









Computational engineering design. Design Automation.

#### Design on demand, Micro Aerial Vehicles, MAV



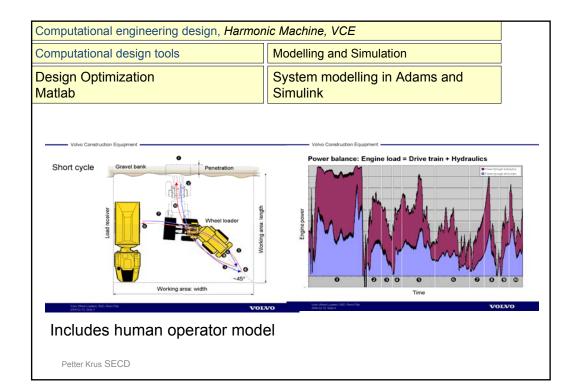
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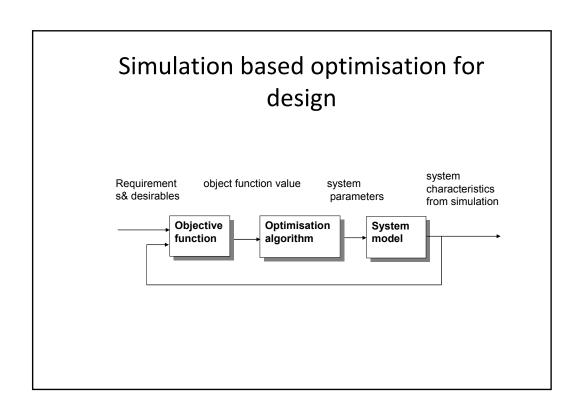
	Baseline	Optimized
motor	Astroflight 010	Hacker B20 31S
propeller	Graupner 6x3 folding	APC 5x5
motor controller	Astroflight 10	YGE 4
Battery	Etec 1200 2s1p	Tanic 1100 3s1p
S	0.12 m2	0.12
AR	1.33	1.33
Endurance at cruise	22 min	42 min
Max speed	75 km/h	
T/W ratio	0.95	1.2

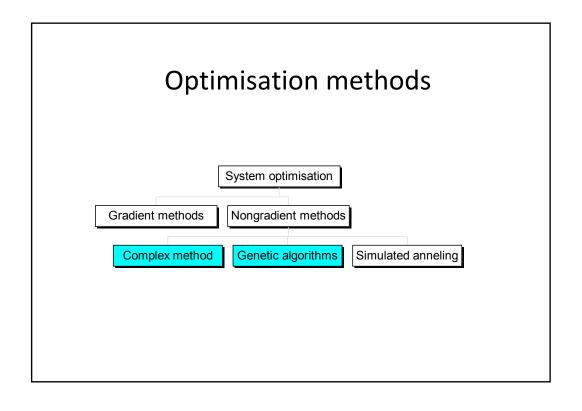


Result: Shopping list and geometrical dimensions

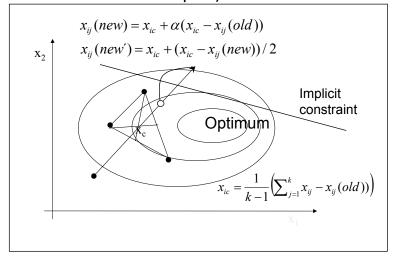
David Lundström, Petter Krus



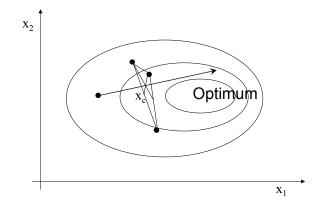


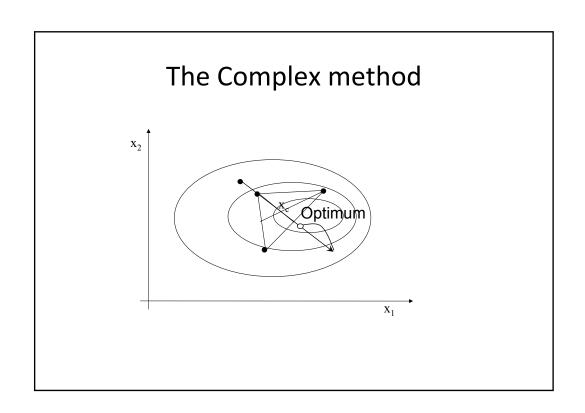


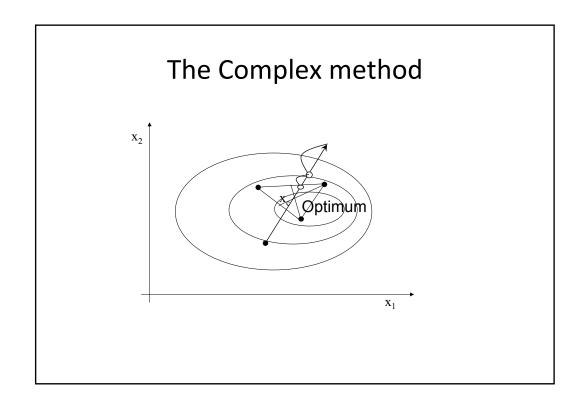
## The Complex method (Constrained simplex)

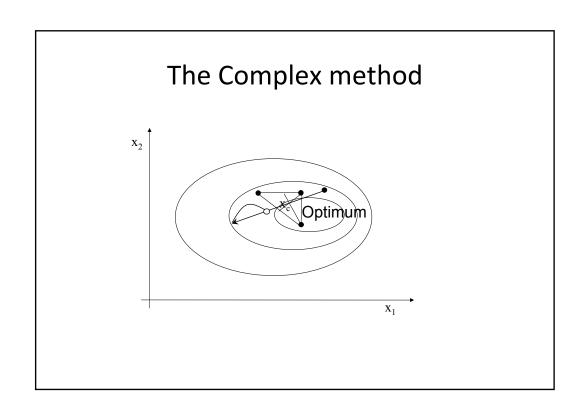


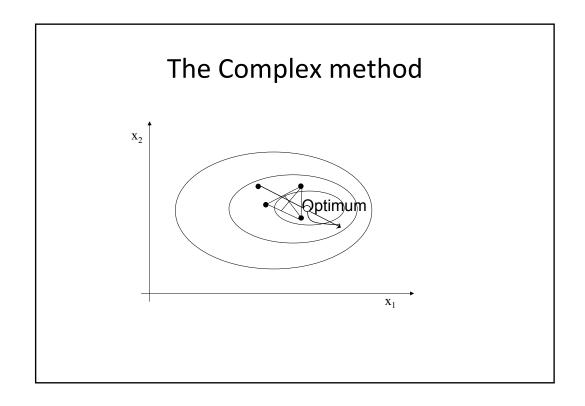




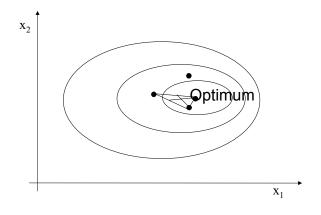




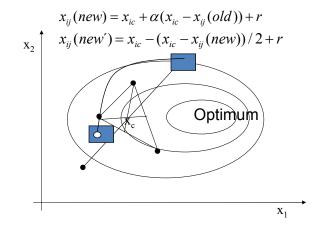








## The Complex-R method



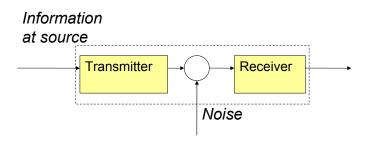
### The Complex-RF

- In order to track dynamic object functions it is necessary to introduce a forgetting factor
- This is done by decreasing an object function value each time it is not re-evaluated
- This is also an advantage in "tricky" nondynamic object functions such as noisy and discrete functions

## Information theory and design

- Both design and design optimization is about increasing information about the product/system
- Information theory was introduced by Shannon in 1947
- Information entropy is a measure of level of uncertainty. High entropy->Low uncertainty

### A general communication system



## Amount of Information content (Information Entropy)

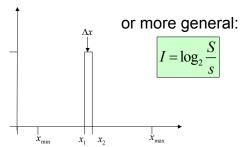
The information content *I* of a variable (in bits).

$$I = \int_{x_{\min}}^{x_{\max}} p(x) \log_2(p(x)x_R) dx \qquad x_R = x_{\max} - x_{\min}$$

If the range  $x_r$  is divided in equal parts  $\Delta x$  the amount of information is:

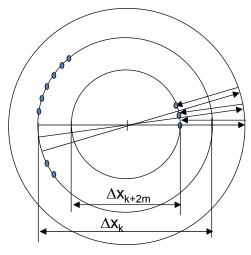
$$I = \log_2 \frac{x_r}{\Delta x} = \log_2 \frac{1}{\delta_x}$$

Here  $\Delta x$  is the tolerance in x.



#### Complex convergence rate

The average degree of contraction in each step



$$\frac{\Delta x_{k+1}}{\Delta x_k} = \left(\frac{\alpha}{2}\right)^{\frac{1}{2m}}$$

$$\frac{\Delta x_{k+1}}{\Delta x_k} = \left(\frac{\alpha}{2}\right)^{\frac{1}{2\kappa n}}$$

$$m = \kappa n$$

## Complex contraction rate

The average degree of contraction in each step of optimization is: 1

$$\frac{\Delta x_{k+1}}{\Delta x_k} = \left(\frac{\alpha}{2}\right)^{\frac{1}{2\kappa n}} \qquad m = \kappa n$$

The increase in information in each step is: (Times n because there are n parameters that are gaining information)

$$\Delta I = -n \log_2 \frac{\Delta x_{k+1}}{\Delta x_k} = -n \log_2 \left(\frac{\alpha}{2}\right)^{\frac{1}{2\kappa n}} = -\log_2 \left(\frac{\alpha}{2}\right)^{\frac{1}{2\kappa}}$$

## **Complex contraction**

The increase in information I in each step

$$\Delta I = -\log_2\left(\frac{\alpha}{2}\right)^{\frac{1}{2\kappa}}$$

Example:  $\alpha$ =1.3,  $\kappa$ =2 yields

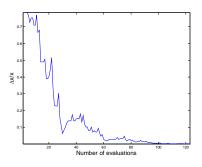
$$\Delta I = -\log_2 \left(\frac{1.3}{2}\right)^{\frac{1}{2*2}} = 0.155$$

The amount of information gained in each step is  $\Delta I = 0.155$  bits/evaluation

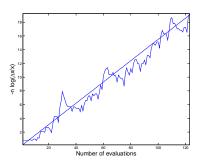
## Test function: Unimodal hump

 $f(x_1, x_2) = \sin(\pi x_1)\sin(\pi x_2)$ 

## Convergence in parameters Unimodal hump



Convergence of optimization variables. This shows the max relative spread dx of all (both) variables



log2 max(dx), and a straight line corresponding to the theoretical convergence rate.

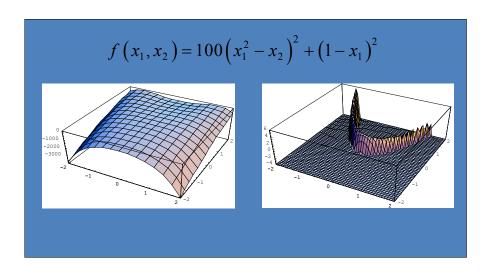
## Efficiency

Unimodal hump

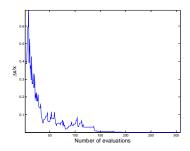
$$\phi = n \frac{\log_2 \varepsilon}{k_m} = -2 \frac{\log_2 0.001}{99} = 0.2$$

bits/evaluation

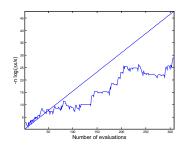
### Test function: Rosenbrock's banana



## Convergence in parameters Banana function



Convergence of optimization variables. This shows the max relative spread dx of all (both) variables



The convergence expressed as – n log2 max(dx), and a straight line corresponding to the theoretical convergence rate.

#### Meta object function.

Minimize the amount of evaluation needed to reach an optimum with a certain probability

$$(1 - P) = (1 - P_{opt})^m$$

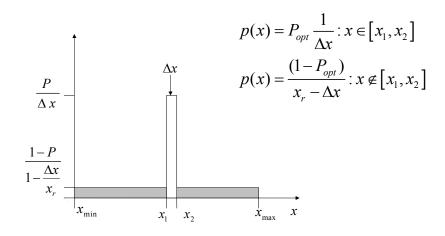
$$m = \frac{Log_2(1-P)}{Log_2(1-P_{out})}$$

 $m = \frac{Log_2(1-P)}{Log_2(1-P_{opt})} \qquad \begin{array}{l} \text{The number of optimization runs} \\ \text{needed to reach optimum with a} \\ \text{certain probability} \end{array}$ 

$$\phi^{(1)} = \frac{I}{mk_m} = -n \frac{\log_2 \varepsilon}{Log_2(1-P)} \frac{Log_2(1-P_{opt})}{k_m}$$
 The adjusted meta object function

$$\widehat{P}_{opt} = hitrate = \frac{N_{succes}}{N}$$

#### Probability distribution after optimization



#### Relative Entropy of the Found Optimum

$$\begin{split} I_{x} &= \int_{x_{\min}}^{x_{\max}} p(x) \log_{2}(p(x)) x_{r} dx = \\ &+ \int_{x_{\min}}^{x_{1}} \frac{(1 - P_{opt})}{1 - \Delta x / x_{r}} \log_{2} \left( \frac{1 - P_{opt}}{1 - \Delta x / x_{r}} \right) dx + \\ &+ \int_{x_{1}}^{x_{2}} P_{opt} \frac{x_{r}}{\Delta x} x \log_{2} \left( P \frac{x_{r}}{\Delta x} \right) dx + \\ &+ \int_{x_{2}}^{x_{\max}} \frac{(1 - P_{opt})}{1 - \Delta x / x_{r}} \log_{2} \left( \frac{(1 - P_{opt})}{1 - \Delta x / x_{r}} \right) dx = \\ &= (1 - P_{opt}) \log_{2} \left( \frac{1 - P_{opt}}{1 - \Delta x / x_{r}} \right) + P_{opt} \log_{2} \left( \frac{P_{opt}}{\Delta x / x_{r}} \right) \end{split}$$

### Meta object function

uncertainty, representing the sum of

 $I_{tot} = (1 - P_{opt}) \log_2 \left( \frac{1 - P_{opt}}{1 - \varepsilon_x^n} \right) + P_{opt} \log_2 \left( \frac{P_{opt}}{\varepsilon_x^n} \right)$ 

uncertainty in location and

I expresses

the total

$$\phi^{(2)} = \frac{I_x}{N_m} = \frac{1}{N_m} \left( (1 - P_{opt}) \log_2 \left( \frac{1 - P_{opt}}{1 - \varepsilon_x^n} \right) + P_{opt} \log_2 \left( \frac{P_{opt}}{\varepsilon_x^n} \right) \right) \approx$$

uncertainty of success

$$(1 - P_{opt}) >> \varepsilon^n$$
  $\phi^{(2)} \approx \frac{1}{N_m} P_{opt} \log_2 \left(\frac{P_{opt}}{\varepsilon_x^n}\right)$ 

### Example: The banana function

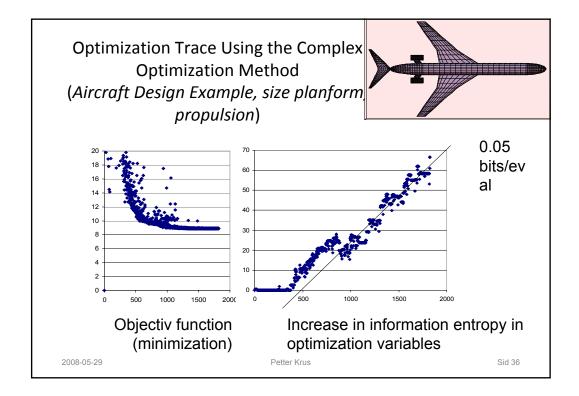
$$\phi^{(2)} = \frac{I_x}{N_{_{AH}}} = \frac{1}{N_{_{AH}}} \left( (1 - P_{_{opt}}) \log_2 \left( \frac{1 - P_{_{opt}}}{1 - \varepsilon_x^{_{A}}} \right) + P_{_{opt}} \log_2 \left( \frac{P_{_{opt}}}{\varepsilon_x^{_{A}}} \right) \right)$$

 $P_{opt} = 0.85$   $\varepsilon = 0.001$  N = 272

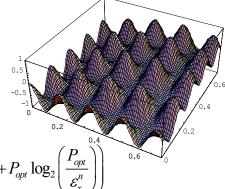
Calculate the efficiency in bits/eval

$$\phi^{(2)} = \frac{I_x}{N_m} = \frac{1}{N_m} \left( (1 - P_{opt}) \log_2 \left( \frac{1 - P_{opt}}{1 - \varepsilon_x^n} \right) + P_{opt} \log_2 \left( \frac{P_{opt}}{\varepsilon_x^n} \right) \right)$$

$$- \frac{1}{272} \left( (1 - 0.85) \log_2 \left( \frac{1 - 0.85}{1 - 0.001^2} \right) + 0.85 \log_2 \left( \frac{0.85}{0.001^2} \right) \right) = 0.044$$
bits/evaluation



## **Example: Multiomodal function**



$$\phi^{(2)} = \frac{I_x}{N_m} = \frac{1}{N_m} \left( (1 - P_{opt}) \log_2 \left( \frac{1 - P_{opt}}{1 - \varepsilon_x^n} \right) + P_{opt} \log_2 \left( \frac{P_{opt}}{\varepsilon_x^n} \right) \right)$$

$$-\frac{2}{133} \left( (1 - 0.78) \log_2 \left( \frac{1 - 0.78}{1 - 0.001^2} \right) + 0.78 \log_2 \left( \frac{0.78}{0.001^2} \right) \right) = 0.0087$$

compare to 0.2 bit/eval for unimodal. 22 times smaller!

bits/evaluation

### Multiomodal function

$$\phi^{(2)} = \frac{I_x}{N_m} = \frac{1}{N_m} \left( (1 - P_{opt}) \log_2 \left( \frac{1 - P_{opt}}{1 - \varepsilon_x^n} \right) + P_{opt} \log_2 \left( \frac{P_{opt}}{\varepsilon_x^n} \right) \right)$$

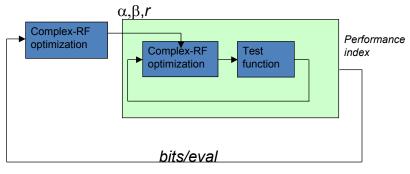
$$- \frac{2}{133} \left( (1 - 0.78) \log_2 \left( \frac{1 - 0.78}{1 - 0.001^2} \right) + 0.78 \log_2 \left( \frac{0.78}{0.001^2} \right) \right) = 0.0087 \text{ bits/evaluation}$$

compare to 0.2 bit/eval for unimodal. 22 times smaller, for 25 times higher modality

• The "curse of dimensionality" is not in the number of parameters. It is in the modality of the output.

### Meta optimization scheme

In each evaluation point optimization of two test functions (Rosenbrock's banana and the integer valued function) where done and the sum of the performance indices where taken as the measure of merit.



## Variations in r and $\gamma$

using Rosenbrock's banana as test function

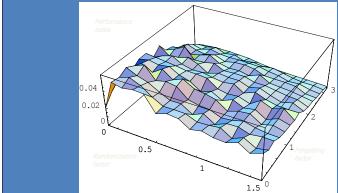


Figure 11. The performance index of the method, as a function of the randomization factor and the forgetting factor for Rosenbrock's banana.

#### Information Entropy

- Design information entropy is a measure of the contraction of design space during design. It is roughly proportional to the number of objects in the design space and thus a more convenient unit than the number of possibilities excluded.
- Relative Information entropy is a very useful state in evolutionary learning processes. It is consistent with the view that the design process is a learning process.
- A general model is presented that links information to optimization, refinement, performance and cost.
- The results indicate that it sometimes may be more efficient to increase the number of design parameters, rather than refining a few parameters to a high degree.

2008-05-29 Petter Krus Sid 41

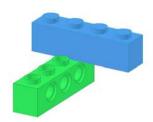
## Information Entropy

- Using the performance criterion the COMPLEX-RF method has been compared with GA and Fmincon(Matlab). The result is surprising in that all the method are remarkably similar in performance index over a wide range of problems (within a factor of three).
- This is in support of the "no free lunch theorem", Wolpert and Macready.

## Design Information Entropy is a Measure of the Size of the Design Space

Lego example

- The design space of a set of Lego bricks represents all combinations of arrangir these bricks.
- With a set of only two bricks with four knobs on each there are 51 discrete possible arrangements
- Two of these represents picking only one brick. And one state is to pick no one.
- The 51 different configuration (states) means that the amount of information needed to specify a particular design is:



$$I_x = \log_2 n_{Dstate} = \log_2 51 = 5.7 \text{bits}$$

2008-05-29 Petter Krus Sid 43

# Amount of Information content (Information Entropy)

The differential information entropy for continuous signals, defined by Shannon [1] as:

$$H = -\int_{-\infty}^{\infty} p(x) \log_2(p(x)) dx$$

Kullback-Leibler divergence

$$H_{rel} = -\int_{-\infty}^{\infty} p(x) \log_2(\frac{p(x)}{m(x)}) dx$$

Generalized

$$H_{rel} = -\int_{-\infty}^{\infty} \cdots \int_{-\infty}^{\infty} p(x_1 \dots x_n) \log_2(\frac{p(x_1 \dots x_n)}{m(x_1 \dots x_n)}) dx_1 \cdots dx_n$$

# Amount of Information content (Information Entropy)

The distribution m(x) should be a rectangular distribution in the bounded interval.

$$I_x = H_{rel}(x) = -\int_{x_{min}}^{x_{max}} p(x) \log_2(p(x)x_R) dx$$

Generalized

$$I_{x} = -\int_{x_{1,\text{min}}}^{x_{1,\text{max}}} \cdots \int_{x_{n,\text{min}}}^{x_{n,\text{max}}} p(x_{1} \dots x_{n}) \log_{2}(p(x_{1}, \dots x_{n}) x_{R1} \cdots x_{Rn}) dx_{1} \cdots dx_{n}$$

More compact

$$I_x = -\int_S p(\mathbf{x}) \log_2(p(\mathbf{x})S) dS$$

2008-05-29 Petter Krus Sid 45

# Amount of Information content (Information Entropy)

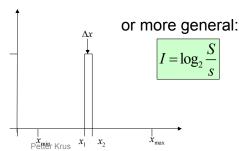
The information content *I* of a variable (in bits).

$$I = -\int_{x_{\min}}^{x_{\max}} p(x) \log_2(p(x)x_R) dx \qquad x_R = x_{\max} - x_{\min}$$

If the range  $x_r$  is divided in equal parts  $\Delta x$  the amount of information is:

$$I = \log_2 \frac{x_r}{\Delta x} = \log_2 \frac{1}{\delta_x}$$

Here  $\Delta x$  is the tolerance in x.



Sid 46

2008-05-29

## Design Space with Both Continuous and Discrete Variables

- The position of the inserted axis represents a continuous variables
- The information entropy associated with that is dependent on the accuracy with which it is specified.

$$I_{x} = \log_{2} n_{Dstates} + \log n_{Cstates} + \log_{2} \frac{x_{R}}{\Delta x}$$

 The axis can be in three position and If the position of the axis within one hole is specified within 10% the total information entropy

$$I_x = \log_2 51 + \log_2 3 + \log_2 \frac{1}{0.1} = 10.6 \text{bits}$$

2008-05-29

Petter Krus

Sid 47

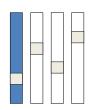
x

## **Design Information entropy**

- The design information entropy represents a measure of the precision by which a design is defined relative to the design space in consideration.
- It is a measure of the dimensionality of the design problem
- It can also be seen as a measure of a probability that a random design in the design space is found, that is within the tolerance region of the design parameters.
- It is hence not possible to know the design information entropy of a design just by looking at it, since it is dependent on the design space from which the design has been picked.
- Design information entropy should not be taken as a direct measure of complexity. A very simple design can represent a large information content if it has been picked from a large design space.

### Design space expansion

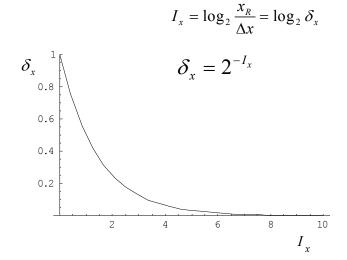
- The design information entropy can be increased in two ways
  - Refinement
  - Design space expansion
- Design space can be increased in several ways like:
  - Adding more bricks
  - Adding other types of bricks
  - Releasing more design parameters in a design



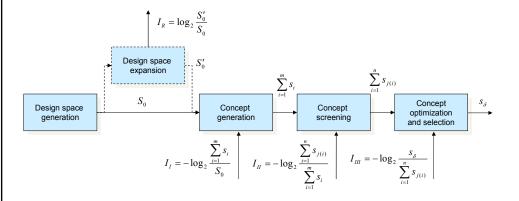
$$I_{x} = \log_{2} \frac{x_{R}'}{\Delta x} = \log_{2} \frac{x_{R}}{\Delta x} \frac{x_{R}'}{x_{R}} = \log_{2} \frac{x_{R}}{\Delta x} \frac{x_{R}^{n}}{x_{R}} = n \log_{2} \frac{x_{R}}{\Delta x}$$
$$= nI_{n}$$

2008-05-29 Petter Krus Sid 49

#### Design uncertainty as a function of design information



## Growth of design information entropy during the design process



2008-05-29 Petter Krus Sid 51

#### Influence factor

$$y = f_x(x)$$
$$z = g(y)$$
$$z = z(x)$$

$$\psi(i) = z_i(x_{R,i}, x_{p,opt,i}) - z_{i-1}(x_{R,i-1}, x_{p,opt,i})$$

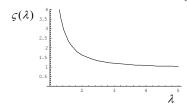
An aproximate modell of this is:

$$\psi(i) = c_p i^{-\lambda}$$

## Asymptotic behaviour

$$z_{opt}(n) = z(n|x_{p,opt}) = \Psi(n) - z_0 =$$
  
=  $c_p \sum_{i=1}^n i^{-k} - z_0 = c_p H(n, \lambda) - z_0$ 

$$\lim_{n \to \infty} H(n, \lambda) = \varsigma(\lambda) = \begin{cases} \infty \text{ when } \lambda \le 1\\ \text{finite value} > 1 \end{cases}$$



2008-05-29 Petter Krus Sid 53

#### Sorted spectrum of parameter influence

$$\psi(i) = c_p i^{-\lambda}$$

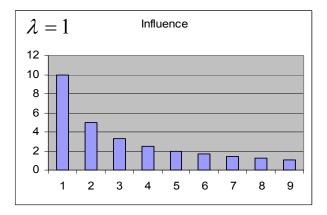


Figure 3. Sorted parameter influences

Total influence up to the n:th parameter:

$$\Psi(n) = \sum_{i=1}^{n} \psi(i) = c_p \sum_{i=1}^{n} i^{-\lambda} = c_p H(n, \lambda)$$

 $H(n,\lambda)$  is the harmonic function

$$z_{opt}(n) = z(n|x_{p,opt}) = \Psi(n) - z_0 =$$

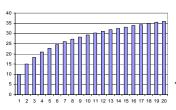
$$= c_p \sum_{i=1}^{n} i^{-k} - z_0 = c_p H(n, \lambda) - z_0$$

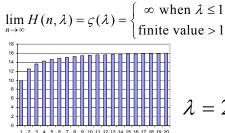
$$\approx c_p \left( c(\lambda) + \frac{1}{\lambda - 1} \left( 1 - \frac{1}{n^{\lambda - 1}} \right) \right) - z_0$$

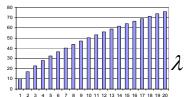
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Sid 55

#### Accumulated influence







#### Performance model

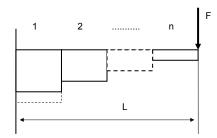
$$p(n) = k_p H(n, \lambda) - p_0$$

$$p_{opt}(n) - p_{opt}(n, \delta_x) = k_p H(n, k) \delta_x \qquad \text{Performance loss} \\ \delta_x = 2^{-\frac{I_x}{n}}$$

$$p_{opt}(n, I_x) = k_p H(n, k)(1 - 2^{\frac{I_x}{n}}) - p_0$$

2008-05-29 Petter Krus Sid 57

#### **Example: optimization of a beam**



$$m = \frac{FL}{\eta k \sigma_{\text{max}}}$$
  $\eta \in ]0,1]$ 

## Efficiency as a function of information

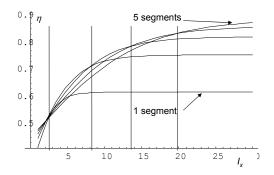
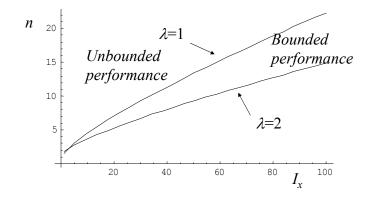
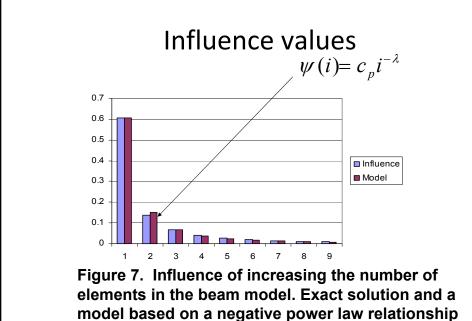


Figure 6. The efficiency of the structure as a function of information  $I_{\rm x}$ 

2008-05-29 Petter Krus Sid 59

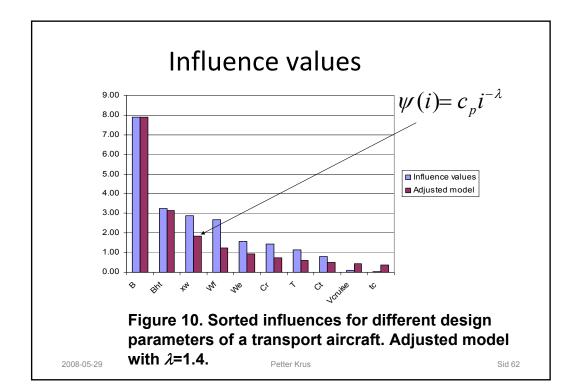
The optimal number of design parameters as a function of design information entropy







Sid 61



### Performance model

$$z_{opt}(n,I_x) = c_p H(n,\lambda)(1-2^{\frac{I_x}{n}}) - z_0$$

Substituting information with cost  $I_x = C/k_c$ 

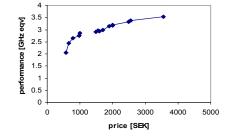
$$z_{opt}(n, C) = c_p H(n, \lambda) (1 - 2^{\frac{C/k_c}{n}}) - z_0$$

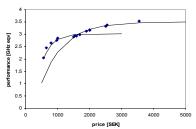
Or as a continuous

$$\begin{aligned} & \text{function} \\ & z_{opt}(n,I_x) = k_p' \, \frac{1}{n^{\lambda-1}} (1-2^{\frac{C}{k_{CI}n}}) - z_0' \end{aligned}$$

2008-05-29 Petter Krus Sid 63

### Example





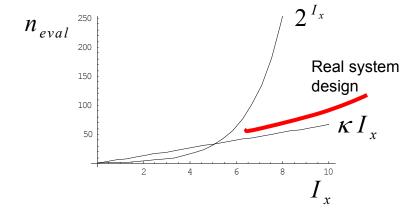
Processor performance as a function of price for Intel Pentium and Celeron processors (2004), with fitted model

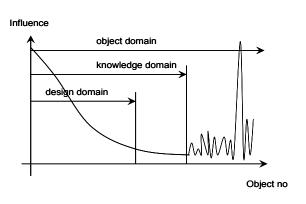
#### Discussion

- Optimality is an illusion of "Flatland", we do not see the constraints that can be removed in the hidden dimmensions.
- The "curse of dimmensionality" is in general resolved by settling for an acceptable solution.
- There is also the "blessing of dimmensionality", many solutions are in fact viable.
- The difference of different optima and suboptima are often shadowed by uncertainty.

2008-05-29 Petter Krus Sid 65

## Number of evaluations needed for a certain amount of information





If previously unknown design parameters enter the knowledge domain the optimal solution might shift into an entirely different optimum, a disruptive design

2008-05-29 Petter Krus Sid 67

## Why is design information entropy useful as a state?

- In the absence of other models, the simplest assumption is that cost is proportional to refinement, which is equivalent to information entropy. This applies both to design cost and manufacturing cost. and should also applicable cost of reliability etc.
- This is supported by the fact that non-gradient optimization methods converge such as the information increases approximately linearly.
- In the absence of more detailed models, the simplest assumption is that performance closes in on its maximum value as the uncertainty goes towards zero.
- It usually have a more convenient magnitude than the number of states.

#### **Neuro Mechanical Neworks**

- There is a general trend for functional density of products to increase.
- Increasing customer requirements
- Prices of microprocessors, sensors and actuators is rapidly decreasing.

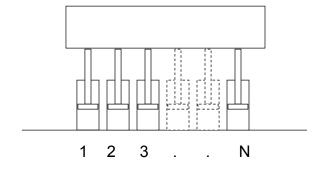
#### Introduction

- 3D-printing techniques
  - Layered manufacturing
- Direct Digital Manufacturing
  - Distributed manufacturing
  - The next big revolution





#### Performance and cost



Performance (force)

$$P = \sum_{i=1}^{n} c_{pi} S_i$$

Cost

$$C = \sum_{i=1}^{n} c_{ci} s_i^{\gamma}$$

If all s are the same

$$P = \sum_{i=1}^{n} c_{p} s = n c_{p} s$$

This yields the component size as:

$$s = \frac{P}{nc_p}$$

In the same way the cost can be calculated

$$C = \sum_{i=1}^{n} c_{c} s^{\gamma} = n c_{c} s^{\gamma} = n c_{c} \left( \frac{P}{n c_{p}} \right)^{\gamma} = n c_{c} \left( \frac{P}{c_{p}} \right)^{\gamma} r^{\gamma} > 1 \text{ drives n to be large and component size to minimum!}$$

$$P = \sum_{i=1}^{n} c_{p} s_{i}^{\gamma_{p}}$$

#### More general

$$C = \sum_{i=1}^{n} c_{c} s_{i}^{\gamma_{c}}$$

If all s are the same

$$P = \sum_{i=1}^{n} c_{p} s = n c_{p} s^{\gamma_{p}} \quad \text{then} \quad s = \left(\frac{P}{n c_{p}}\right)^{\frac{1}{\gamma^{p}}}$$

In the same way the cost can be calculated/

$$C = \sum_{i=1}^{n} c_{c} s^{\gamma_{c}} = n c_{c} s^{\gamma_{c}} = n c_{c} \left(\frac{P}{n c_{p}}\right)^{\frac{\gamma_{c}}{\gamma_{p}}} = n^{\frac{1-\gamma_{c}}{\gamma_{p}}} c_{c} \left(\frac{P}{c_{p}}\right)^{\frac{\gamma_{c}}{\gamma_{p}}}$$

$$component size$$
to minimum!

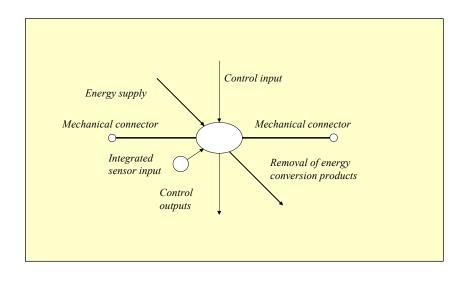
#### Consequences

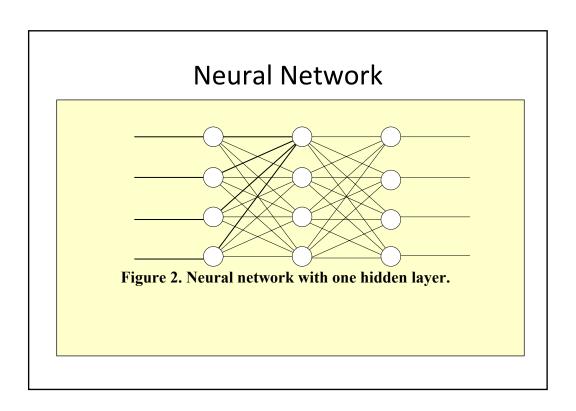
- In electronics  $\gamma_c$  is already greater than  $\gamma_p$  since  $\gamma_p$  is close to zero. Performance is much more related to number than size.
- As this happens in other technologies, component size will be driven down to the minimum size.
- This has happened in biological systems
- This leads to hyper functional technical products

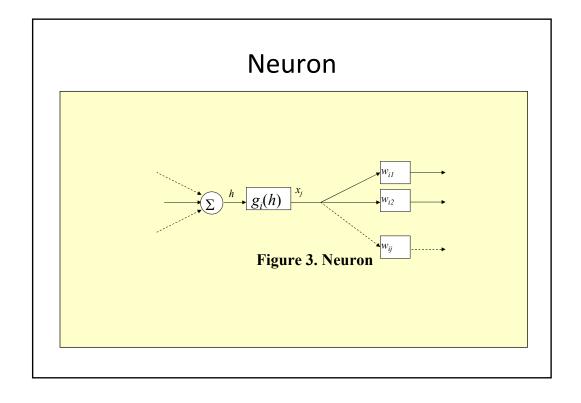
## The Generic Integrated Actuator

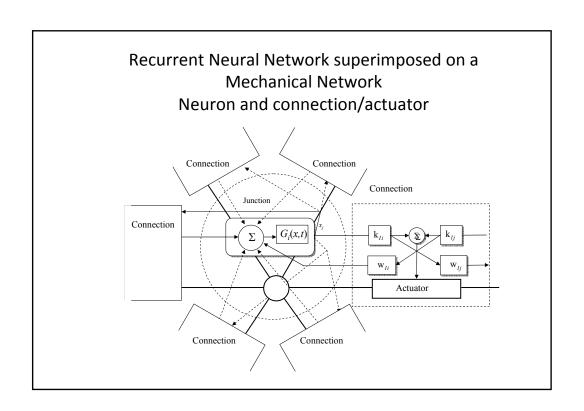
- Actuation, through mechanical connections
- Energy supply and removal of energy conversion residues
- Control and/or sensory input
- Sensory output
- · Simple signal processing

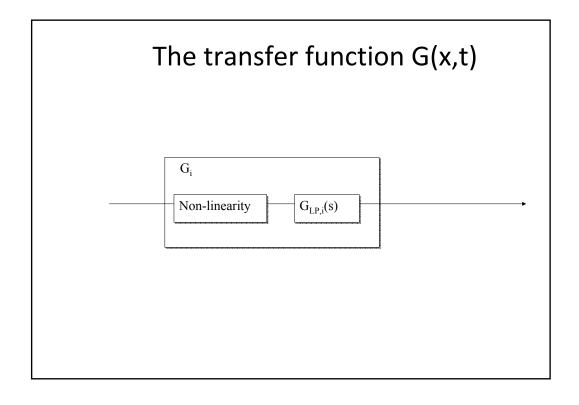
## The Generic Integrated Actuator











### The neuro-mechanical network

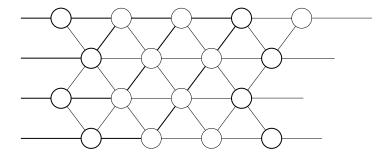
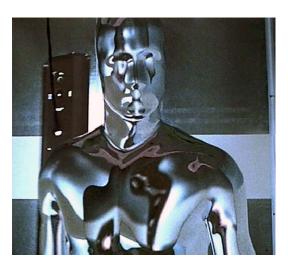


Figure 4 Neuro-mechanical network corresponding to the first two layers in the neural network.

## **Applications**

- As a generic structure using optimisation for mechatronic system design
- As a model for future massively redundant mechatronic systems
- As an explanation model for biological systems

## Application?



## Example

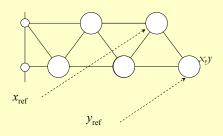


Figure 7. A neuro-mechanical network with five neuron, eight actuators and two inputs.

## The objective function

$$f_{obj} = \int (x - x_{ref})^2 dt + \int (y - y_{ref})^2 dt$$

All the 30 static gains are optimisation variables

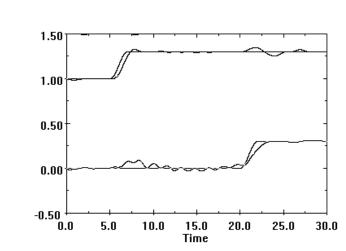
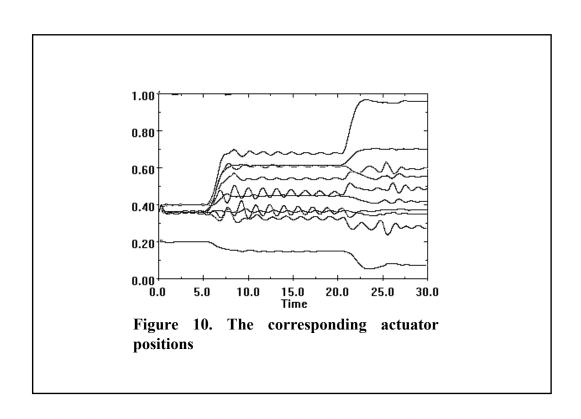
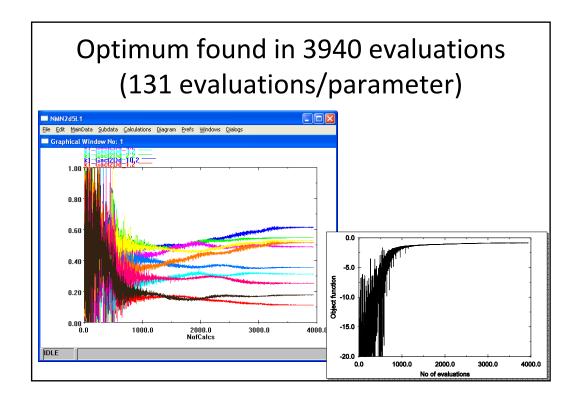


Figure 9. The reference and actual positions in x-and y-directions at the tip.





#### Discussion

- The Neuro mechanical elements can be viewed as engineering stem cells
- They can be used to produce "organic" structures, and parts can have degrees of freedom to be specialised as signal processing, actuators, and pure structure.
- They can also be used to generate more conventional systems through topology optimisation

## Publications related to the topics

- Krus P. 'The Relation Between Performance, Design Space, Refinement and Information'. International Design Engineering Technical Conferences Computers and Information in Engineering Conference ASME/IDETC, Las Vegas, Nevada, USA, 2007.
- Krus P., Andersson J. "An Information Theoretical Perspective on Design Optimization" Proceedings of 2004 DETC:Design Engineering Technical Conference. Salt Lake city, Utah.
- Krus P, Karlsson M,'Neuro-Mechanical Networks Self-Organizing Multifunctional Systems' Presented at the 'Ninth Bath International Fluid Power Workshop', Bath, UK 2002.